POLITECNICO DI TORINO

(01JEUHT) Formal Languages and Compilers

Laboratory N° 2

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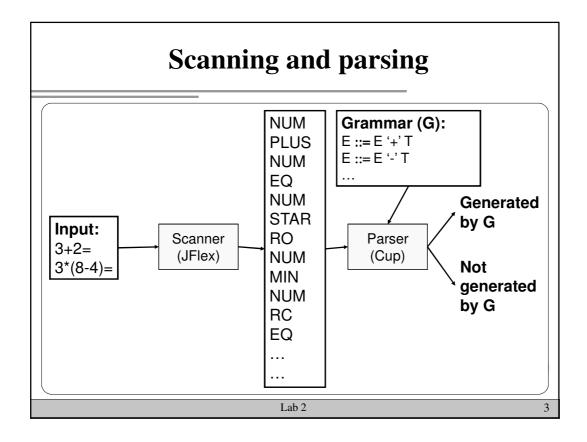


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Parser and syntax analyzer

- Given a non-ambiguous grammar and a sequence of input symbols, a parser is a program that verifies whether the sequence can be generated by means of a derivation from the grammar.
- A syntax analyzer (parser) is a program capable of associating to the input sequence the correct parse tree.
- Parsers can be classified as
 - top-down (parse tree is built from the root to the leaves)
 - bottom-up (parse tree is built from the leaves to the root) : CUP





Context-Free Grammar Definition

A CF grammar is described by

- T, NT, S, PR
- T: Terminals / tokens of the language
- NT: Non-terminals
 - Denote sets of strings generated by the grammar
- S: Start symbol
 - $S \in NT$
- PR: Production rules
 - Indicate how T and NT are combined to generate valid strings
 - PR: NT ::= T | NT



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Example

Derivation:

■ A sequence of grammar rule applications and substitutions that transform a starting non-terminal into a sequence of terminals (tokens).

```
assign_stmt ::= ID EQ expr S;
expr ::= expr operator term;
expr ::= term;
term ::= ID;
term ::= FLOAT;
term ::= INT;
operator ::= PLUS;
operator ::= MIN;
```



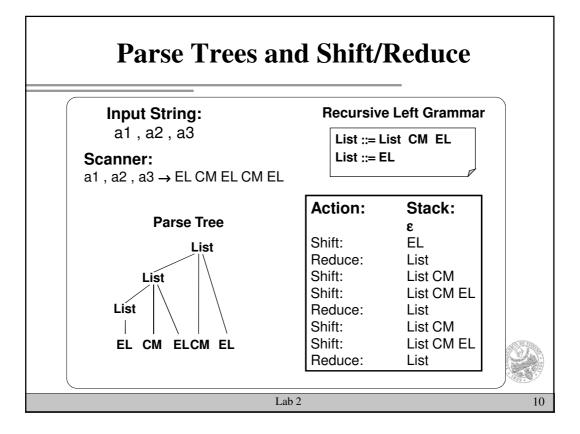
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How bottom-up parsing works: Shift/Reduce tecnique

- A stack, initially empty, is used to keep track of symbols already recognized.
- Terminal symbols are pushed in the stack (shift), until the top of the stack contains a handle (right hand side of a production): the handle is then substituted by the corresponding nonterminal (reduce).
- Note that the reduce operation may only be applied to the top of the stack.
- Parsing is successful only when at the end of the input stream the stack contains only the start symbol





Introduction to CUP

- Cup is a parser generator that transforms the definition of a context-free grammar in a Java program that parses sequences of input symbols according to the grammar itself.
- Besides defining syntax rules, it is possible to specify actions to be executed whenever a production is reduced.
- The parser must be integrated by a scanner: some conventions simplify the integration of Cup-generated parses with JFlexgenerated scanners.
- Official manual: http://www2.cs.tum.edu/projects/cup/



Source file format

- A Cup source file has a syntax very similar to Java programs.
- It can be ideally divided in the following sections:
 - Setup
 - Terminals and non-Terminals
 - Precedences (Next lesson)
 - Rules
- Comments are allowed following Java syntax (included in /* and */, or preceded by //)



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Setup section

• This section contains all the directives needed for the parser

- Inclusion of Cup library and other libraries: import java_cup.runtime.*;
- User code: (Next lesson)
 - Ridefinition of Cup internal methods
 - Integration with scanner other than JFlex



Terminals / Non-Terminals section

- It contains the definition of
 - Terminals: passed by JFlex
 - Non-Terminals
 - The grammar start symbol
- Start symbol
 - start with <non_terminal_name>;
 - It is the root of the parse tree
 - Only one occurrence of this keyword is allowed



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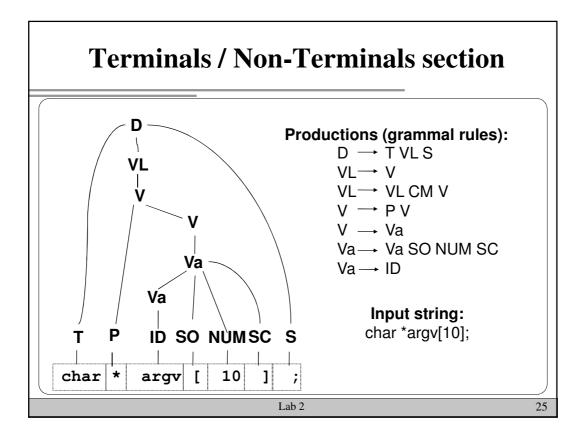
Terminals / Non-Terminals section

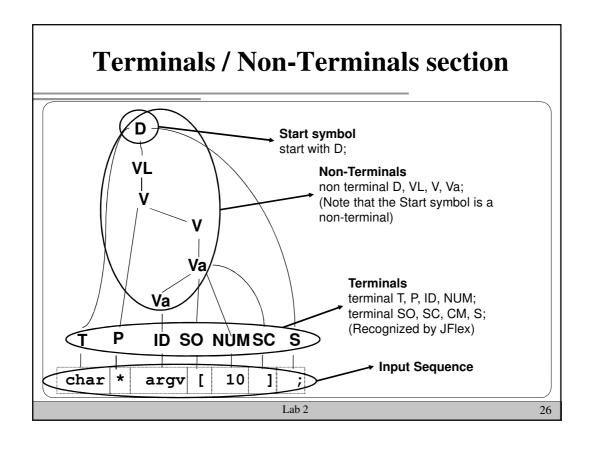
Terminals

- terminal <terminal_1>,...,<terminal_n> ;
 - <terminal>: name containing letters, '_', '.' and digits (the first character must be a letter)
- Terminals are recognized by Jflex
- Non-Terminals
 - non terminal <non_terminal_1>,...,<non_terminal_n> ;
 - <non_terminal>: name containing letters, '_', '.' and digits (the first character must be a letter).



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Rules section

• The Rules section contains one or more productions in the form:

```
<non_terminal> ::= Right_Hand_Side ;
```

- where Right_Hand_Side is a sequence of 0 or more symbols.
- To each prodution, an action can be associated, which must be enclosed between {: and :}
 - Note: the action is executed just before the reduce operation takes place
- Example:

```
D ::= T VL S
     {: System.out.println("Declaration found"); :}
;
```



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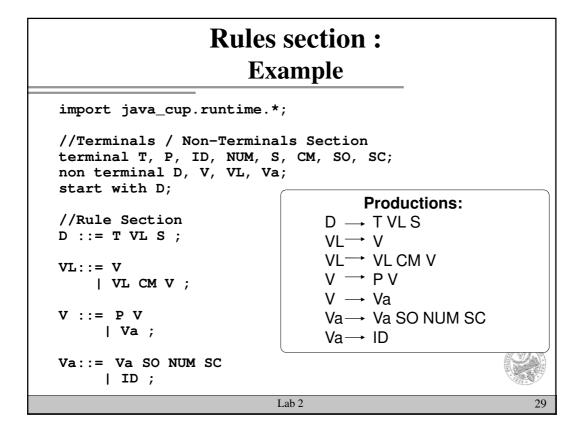
Rules section (2)

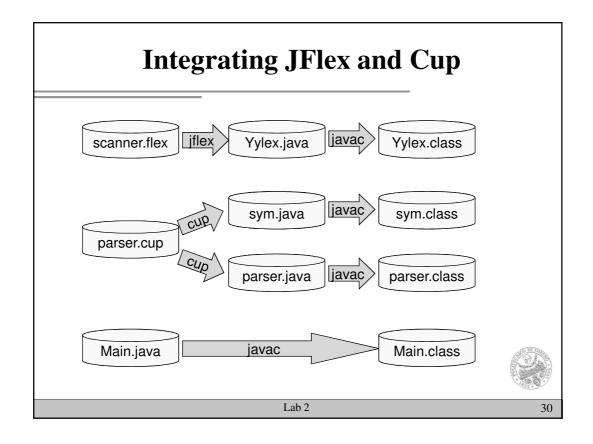
- If more than one production exist for a given non-terminal, they must be grouped and separated by '|'.
- Es.

■ NB: the use of the "|" character generates two separates rules. It is important to remember that the code between {: and :} is executed only when a giver rule is matched.

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Integrating JFlex and Cup

- Parser and scanner must agree on the values associated to each token (terminal)
- When the scanner recognizes a token, it must pass a suitable value to the parser. This is done by means of the Symbol class, whose constructors are:
 - public Symbol(int sym id)
 - public Symbol(int sym_id, int left, int right)
 - public Symbol(int sym_id, Object o)
 - public Symbol(int sym id, int left, int right, Object o)
 - The class Symbol can be found in the cup installation directory:
 - Java_cup/runtime/Symbol.java
- When a terminal is defined by means of the terminal keyword, Cup associated an integer value to that token.
 - This mapping is contained in the file sym.java generated by cup during the compiling process

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Integrating JFlex and Cup (2)

• If in the parser the following list of terminal symbols has been declared:

```
terminal T, P, ID, NUM, PV, CM, SO, SC, S;
```

They can be used inside the scanner and passed to the parser in the following way:

```
[a-zA-Z_{-}][a-zA-Z0-9_{-}]*
                                 {return new Symbol(sym.ID);}
1/
                                 {return new Symbol(sym.SO);}
                                 {return new Symbol(sym.SC);}
/]
"/*" ~ "*/"
                                 {;}
                                 {;}
\r | \n | \r\n | " " | \t
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```

