


POLITECNICO DI TORINO

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**(01JEUHT) Formal Languages and Compilers**  
Laboratory N° 2

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Lab 2 1

## Parser and syntax analyzer

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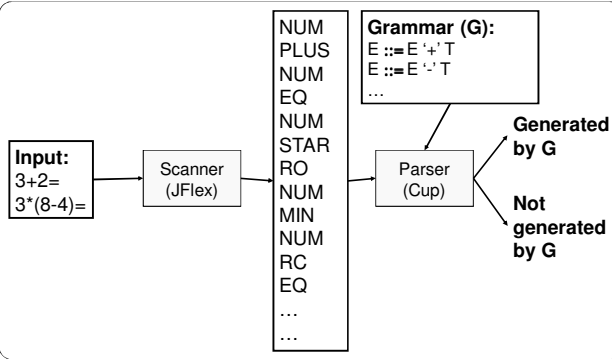
- Given a non-ambiguous grammar and a sequence of input symbols, a parser is a program that verifies whether the sequence can be generated by means of a derivation from the grammar.
- A syntax analyzer (parser) is a program capable of associating to the input sequence the correct parse tree.
- Parsers can be classified as
  - top-down (parse tree is built from the root to the leaves )
  - bottom-up (parse tree is built from the leaves to the root ) : CUP



Lab 2 2

## Scanning and parsing

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Lab 2 3

## Context-Free Grammar Definition

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A CF grammar is described by

- T, NT, S, PR
- T: Terminals / tokens of the language
- NT: Non-terminals
  - Denote sets of strings generated by the grammar
- S: Start symbol
  - $S \in NT$
- PR: Production rules
  - Indicate how T and NT are combined to generate valid strings
    - PR:  $NT ::= T \mid NT$



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## Example

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- **Derivation:**
  - A sequence of grammar rule applications and substitutions that transform a starting non-terminal into a sequence of terminals (tokens).

```

assign_stmt ::= ID EQ expr S ;
expr ::= expr operator term ;
expr ::= term ;
term ::= ID ;
term ::= FLOAT ;
term ::= INT ;
operator ::= PLUS ;
operator ::= MIN ;
    
```




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## How bottom-up parsing works: Shift/Reduce technique

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- A stack, initially empty, is used to keep track of symbols already recognized.
- Terminal symbols are pushed in the stack (shift), until the top of the stack contains a handle (right hand side of a production): the handle is then substituted by the corresponding non-terminal (reduce).
- Note that the reduce operation may only be applied to the top of the stack.
- Parsing is successful only when at the end of the input stream the stack contains only the start symbol



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### Parse Trees and Shift/Reduce

**Input String:**  
a1 , a2 , a3

**Scanner:**  
a1 , a2 , a3 → EL CM EL CM EL

**Recursive Left Grammar**

```
List ::= List CM EL
List ::= EL
```


**Action: Stack:**

Shift:	ε
Reduce:	EL
Shift:	List
Shift:	List CM
Shift:	List CM EL
Reduce:	List
Shift:	List CM
Shift:	List CM EL
Reduce:	List

**Parse Tree**


```

      List
     /  |  \
    List |  |
   /  |  |  \
  List |  |  |  \
 EL CM |  |  |  \
      |  |  |  \
      EL CM EL
    
```




### Introduction to CUP

- Cup is a parser generator that transforms the definition of a context-free grammar in a Java program that parses sequences of input symbols according to the grammar itself.
- Besides defining syntax rules, it is possible to specify actions to be executed whenever a production is reduced.
- The parser must be integrated by a scanner: some conventions simplify the integration of Cup-generated parses with JFlex-generated scanners.
- Official manual:  
<http://www2.cs.tum.edu/projects/cup/>




### Source file format

- A Cup source file has a syntax very similar to Java programs.
- It can be ideally divided in the following sections:
  - Setup
  - Terminals and non-Terminals
  - Precedences (Next lesson)
  - Rules
- Comments are allowed following Java syntax (included in /\* and \*/ , or preceded by //)




### Setup section

- This section contains all the directives needed for the parser
- Inclusion of Cup library and other libraries:  
import java\_cup.runtime.\*;
- User code: (Next lesson)
  - Rdefinition of Cup internal methods
  - Integration with scanner other than JFlex




### Terminals / Non-Terminals section

- It contains the definition of
  - Terminals: passed by JFlex
  - Non-Terminals
  - The grammar start symbol
- Start symbol
  - start with <non\_terminal\_name> ;
  - It is the root of the parse tree
  - Only one occurrence of this keyword is allowed



### Terminals / Non-Terminals section

- Terminals
  - terminal <terminal\_1>,...,<terminal\_n> ;
    - <terminal>: name containing letters, '\_', '.' and digits (the first character must be a letter)
  - Terminals are recognized by Jflex
- Non-Terminals
  - non terminal <non\_terminal\_1>,...,<non\_terminal\_n> ;
    - <non\_terminals>: name containing letters, '\_', '.' and digits (the first character must be a letter).



### Terminals / Non-Terminals section

**Productions (grammal rules):**

$D \rightarrow T VL S$   
 $VL \rightarrow V$   
 $VL \rightarrow VL CM V$   
 $V \rightarrow P V$   
 $V \rightarrow Va$   
 $Va \rightarrow Va SO NUM SC$   
 $Va \rightarrow ID$

**Input string:**  
char \*argv[10];

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### Terminals / Non-Terminals section

**Productions (grammal rules):**

$D \rightarrow T VL S$   
 $VL \rightarrow V$   
 $VL \rightarrow VL CM V$   
 $V \rightarrow P V$   
 $V \rightarrow Va$   
 $Va \rightarrow Va SO NUM SC$   
 $Va \rightarrow ID$

**Input string:**  
char \*argv[10];

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### Rules section

- The Rules section contains one or more productions in the form:
 

```
<non_terminal> ::= Right_Hand_Side ;
```
- where *Right\_Hand\_Side* is a sequence of 0 or more symbols.
- To each production, an action can be associated, which must be enclosed between { : and ; }
  - Note: the action is executed just before the reduce operation takes place
- Example:
 

```
D ::= T VL S
           { System.out.println("Declaration found"); ; }
```

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### Rules section (2)

- If more than one production exist for a given non-terminal, they must be grouped and separated by '|'.
 

```
funz ::= type ID RO VL RC S
           { System.out.println("Function prototype"); ; }
           | type ID RO VL RC BO stmt_list BC
           { System.out.println("Function"); ; }
```
- NB: the use of the "|" character generates two separates rules. It is important to remember that the code between { : and ; } is executed only when a giver rule is matched.

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### Rules section : Example

```
import java_cup.runtime.*;

//Terminals / Non-Terminals Section
terminal T, P, ID, NUM, S, CM, SO, SC;
non terminal D, V, VL, Va;
start with D;

//Rule Section
D ::= T VL S ;

VL ::= v
      | VL CM V ;

V ::= P V
     | Va ;

Va ::= Va SO NUM SC
     | ID ;
```

**Productions:**

$D \rightarrow T VL S$   
 $VL \rightarrow V$   
 $VL \rightarrow VL CM V$   
 $V \rightarrow P V$   
 $V \rightarrow Va$   
 $Va \rightarrow Va SO NUM SC$   
 $Va \rightarrow ID$

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### Integrating JFlex and Cup

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## Integrating JFlex and Cup

- Parser and scanner must agree on the values associated to each token (terminal)
- When the scanner recognizes a token, it must pass a suitable value to the parser. This is done by means of the Symbol class, whose constructors are:
  - public Symbol( int sym\_id)
  - public Symbol( int sym\_id, int left, int right)
  - public Symbol( int sym\_id, Object o)
  - public Symbol( int sym\_id, int left, int right, Object o)
  - The class Symbol can be found in the cup installation directory:
    - Java\_cup/runtime/Symbol.java
- When a terminal is defined by means of the terminal keyword, Cup associated an integer value to that token.
  - This mapping is contained in the file sym.java generated by cup during the compiling process

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## Integrating JFlex and Cup (2)

- If in the parser the following list of terminal symbols has been declared:
 

```
terminal T, P, ID, NUM, PV, CM, SO, SC, S;
```
- They can be used inside the scanner and passed to the parser in the following way:

```
...
%%
...
%%
[a-zA-Z_][a-zA-Z0-9_]* {return new Symbol(sym.ID);}
\[ {return new Symbol(sym.SO);}
\\ {return new Symbol(sym.SC);}
...
"/" ~ "*" / {;}
\r | \n | \r\n | " " | \t {;}
```

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scanner.flex

## Scanner modifications

- Include the Cup library ( java\_cup.runtime.\* ) in the code section
- Activate Cup compatibility by means of the %cup directive in the Declarations section

```
import java_cup.runtime.*;
...
%%
%cup
...
%%
[a-z]+ { return new Symbol(sym.EL); }
"," { return new Symbol(sym.CM); }
```

List → List CM EL  
List → EL

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parser.cup

## The Cup parser

```
import java_cup.runtime.*;
```

```
terminal EL, CM;
non terminal List, EList;
```

List → List CM EL  
List → EL

```
start with EList;
```

```
EList ::= List {; System.out.println("List found"); ;} |
        {; System.out.println("Empty list"); ;}
```

```
List ::= List CM EL
```

```
List ::= EL
```

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Main.java

## Main

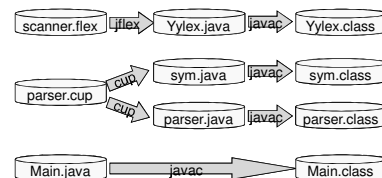
```
import java.io.*;

public class Main {
    static public void main(String argv[]) {
        try {
            /* Instantiate the scanner and open input file argv[0] */
            Yylex l = new Yylex(new FileReader(argv[0]));
            /* Instantiate the parser */
            parser p = new parser(l);
            /* Start the parser */
            Object result = p.parse();
        } catch (Exception e) {
            e.printStackTrace();
        }
    }
}
```

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## Compiling



- jflex scanner.jflex
- java java\_cup.Main parser.cup
  - In the case of shift/reduce or reduce/reduce conflicts:
    - java java\_cup.Main -expect <number\_of\_conflicts> parser.cup
    - java java\_cup.MainDrawTree parser.cup
      - Can be used in LABINF or at home installing a modified version of the parser
      - The parse tree is drawn (useful for debugging)

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