

POLITECNICO DI TORINO

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## (01JEUHT) Formal Languages and Compilers

### Laboratory N°4

Stefano Scanzio  
mail: [stefano.scanzio@polito.it](mailto:stefano.scanzio@polito.it)  
Web: <https://www.skenz.it/compilers>



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## Attributes of Symbols

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- A set of attributes can be associated to each symbol; attributes can be:
  - **Synthesized**: calculated from the values of the attributes of the node's children in the parse tree,
  - **Inherited**: calculated from the values of the parents / siblings in the parse tree.
- A set of semantic rules, specifying how attributes are calculated, is associated to each production.
- The scanner passes semantic values to the parser which, while recognizing the grammar, updates the nodes of the parse tree



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## Synthesized attributes

- A grammar whose attributes are all synthesized is denoted as an S-attribute grammar.
- In this case, it is possible to calculate the values of all attributes using a bottom-up strategy, from the leaves to the root of the parse tree.

$E ::= E_1 '+' T$	$E.value = E_1.value + T.value$
$E ::= T$	$E.value = T.value$
$T ::= number$	$T.value = number.value$



## Cup & Semantics: the Symbol class

- In Cup, each symbol in the stack is an object of class Symbol ( cup/java\_cup/runtime/Symbol.java )
- It contains the following information:
  - A number uniquely identifying the symbol
    - ↳ public int sym;
  - The state in which the parse is
    - ↳ public int parse\_state;
  - Two integers that are used to pass the line and column number from the scanner to the parser
    - ↳ public int left, right;
  - An object of class Object to handle semantics
    - ↳ public Object value;



scanner.jflex

## Passing semantic values to the parser

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- Symbol and semantic value:
 


```
[a-zA-Z][a-zA-Z0-9_]* { return new Symbol(sym.ID, new String(yytext())); }
```
- Symbol, line number, column number, and semantic value:
 

```
%{
  private Symbol my_symbol(int type, Object value){
    return new Symbol(type, yyline, yycolumn, value);
  }
}%
[a-zA-Z][a-zA-Z0-9_]* { return my_symbol(sym.ID, new String(yytext())); }
```

**Symbol Constructors:**

```
public Symbol( int sym_id)
public Symbol( int sym_id, int left, int right)
public Symbol( int sym_id, Object o)
public Symbol( int sym_id, int left, int right, Object o)
```
- Or equivalently:
 

```
[a-zA-Z][a-zA-Z0-9_]* {
  return new Symbol(sym.ID, yyline, yycolumn, new String(yytext())); }
```




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## Cup & Semantic: specifying nodes types

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- Cup must know the type of the semantic value of each symbol
- It uses the following definition of terminals and non-terminals:
  - terminal **<Object>** <list\_of\_terminals> ;
  - non terminal **<Object>** <list\_of\_not\_terminals> ;
- **<Object>** is the class of the object associated to a given symbol
- Example:
  - terminal String ID;
    - ↗ An object of class String will be associated to ID.
  - terminal Integer NUM;
  - non terminal MyObject var;

```
class MyObject {
  public String var_name;
  public String var_type;
}
```



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## Cup & Semantic: using semantic values

- Given a set of productions:

```
E ::= E PLUS T
    | E MINUS T ;
```

- One can refer to the semantic value of each symbol by adding labels to the symbols of interest:

- A label is constituted by the ':' character followed by a name

```
E ::= E:n1 PLUS T:n2
    | E:n1 MINUS T:n2 ;
```

- Within each production, the labels can be used normally as objects of the class specified in the definition of terminals and non-terminals:

```
E ::= E:n1 PLUS T:n2 { : System.out.print(n1 + " + " + n2); ; }
    | E:n1 MINUS T:n2 { : System.out.print(n1 + " - " + n2); ; }
```



## Cup & Semantic: Actions and RESULT

- An action can be associated to each production, ( `{: /* Java Code*/ :}` ) and is executed every time the corresponding production is reduced
- The action updates the semantic value of each symbol
- For each production, the RESULT object, of class Object, is defined.
- RESULT represents the result of the semantic rules contained in the action, **and is therefore associated to the symbol in the left hand side of the production**



## Calculating synthesized attributes

- Given the algebraic expressions grammar, the following rule assigns to the symbol 'E' the sum or the subtraction of the values of the addends/subtrahends:

```

non terminal Integer E, T;

E ::=  E:n1 PLUS T:n2
      | E:n1 MINUS T:n2
      ;
      { : RESULT = n1 + n2; : }
      { : RESULT = n1 - n2; : }

```

- OR: { : RESULT = new Integer(n1.intValue() + n2.intValue()); : }



## Calculating synthesized attributes (2)

- It is possible to propagate more than one semantic value through RESULT, in the following way:

```

terminal RO, RC;
terminal String identifier;
terminal Integer Args;
non terminal Object[ ] Func;
non terminal goal;

goal ::= Func:a {
    System.out.println("Function name: " + a[0] + "Number of parameters: " + a[1]);
};

Func ::= identifier:a RO Args:b RC {
    RESULT = new Object[2];
    RESULT[0] = new String(a);
    RESULT[1] = new Integer(b);
};

```



## Calculating synthesized attributes (3)

- Alternatively, one can write a class that contains all the required information:

```

action code {
  class MyFunc {
    public String id;
    public Integer args;
    MyFunc(String id, Integer args) {
      this.id = new String(id);
      this.args = new Integer(args);
    }
  }

  non terminal MyFunc Func;

  goal ::= Func:a {
    System.out.println( "Function name : " + a.id + "Number of parameters: " + a.args );
  };

  Func ::= identifier:a RO Args:b RC { RESULT = new MyFunc( a, b ); };

```

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## Parser debugging

- A series of option are available in Cup to visualize the parser's internal structures:
  - -dump\_grammar : Prints the list of terminals, non-terminals and productions
  - -dump\_states : Prints the state graph
  - -dump\_table : Prints the ACTION TABLE and the REDUCE TABLE
  - -dump : Prints all information
- The parser can be executed in debug mode (all the actions performed to analyze the input sequence are printed)

### Normal mode:

```

Yylex l = new Yylex(new FileReader(file));
parser p = new parser(l);
Object result = p.parse();

```

### Debug mode:

```

Yylex l = new Yylex(new FileReader(file));
parser p = new parser(l);
Object result = p.debug_parse();

```



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## States

● **-dump\_states**

```

===== Viable Prefix Recognizer =====
START lalr_state [0]: {
  [exp ::= (*) T , {EOF PLUS }]
  [exp ::= (*) exp PLUS T , {EOF PLUS }]
}
[T ::= (*) NUMBER , {EOF PLUS }]
[$START ::= (*) exp EOF , {EOF }]
}
transition on exp to state [3]
transition on T to state [2]
transition on NUMBER to state [1]
-----
lalr_state [1]: {
  [T ::= NUMBER (*) , {EOF PLUS }]
}
-----
lalr_state [2]: {
  [exp ::= T (*) , {EOF PLUS }]
}
-----
lalr_state [3]: {
  [exp ::= exp (*) PLUS T , {EOF PLUS }]
  [$START ::= exp (*) EOF , {EOF }]
}
transition on EOF to state [5]
transition on PLUS to state [4]
-----
lalr_state [4]: {
  [exp ::= exp PLUS (*) T , {EOF PLUS }]
  [T ::= (*) NUMBER , {EOF PLUS }]
}
transition on T to state [6]
transition on NUMBER to state [1]
-----
lalr_state [5]: {
  [$START ::= exp EOF (*) , {EOF }]
}
-----
lalr_state [6]: {
  [exp ::= exp PLUS T (*) , {EOF PLUS }]
}
-----
                    
```

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## Action / Reduce Tables

● **-dump\_tables**

```

----- ACTION_TABLE -----
From state #0
 [term 2:SHIFT(to state 1)]
From state #1
 [term 0:REDUCE(with prod 3)] [term 3:REDUCE(with prod 3)]
From state #2
 [term 0:REDUCE(with prod 2)] [term 3:REDUCE(with prod 2)]
From state #3
 [term 0:SHIFT(to state 5)] [term 3:SHIFT(to state 4)]
From state #4
 [term 2:SHIFT(to state 1)]
From state #5
 [term 0:REDUCE(with prod 0)]
From state #6
 [term 0:REDUCE(with prod 1)] [term 3:REDUCE(with prod 1)]
-----
                    
```

```

----- REDUCE_TABLE -----
From state #0
 [non term 1->state 3] [non term 2->state 2]
From state #1
From state #2
From state #3
From state #4
 [non term 2->state 6]
From state #5
From state #6
-----
                    
```

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## Grammar

### ● -dump\_grammar

```

===== Terminals =====
[0]EOF [1]error [2]NUMBER [3]PLUS
===== Non terminals =====
[0]$START [1]exp [2]T
===== Productions =====
[0] $START ::= exp EOF
[1] exp ::= exp PLUS T
[2] exp ::= T
[3] T ::= NUMBER

```

```

exp → exp PLUS T
exp → T
T → NUMBER

```



## Debugging

### ● debug\_parse()

```

# Initializing parser
FOUND: 3
# Current Symbol is #2
# Shift under term #2 to state #1
FOUND: +
# Current token is #3
# Reduce with prod #3 [NT=2, SZ=1]
# Reduce rule: top state 0, lhs sym 2 -> state 2
# Goto state #2
# Reduce with prod #2 [NT=1, SZ=1]
# Reduce rule: top state 0, lhs sym 1 -> state 3
# Goto state #3
# Shift under term #3 to state #4
FOUND: 5
# Current token is #2

```

### Input string: 3+5

```

# Shift under term #2 to state #1
# Current token is #0
# Reduce with prod #3 [NT=2, SZ=1]
# Reduce rule: top state 4, lhs sym 2 -> state 6
# Goto state #6
Found expression
# Reduce with prod #1 [NT=1, SZ=3]
# Reduce rule: top state 0, lhs sym 1 -> state 3
# Goto state #3
-----
# Shift under term #0 to state #5
# Current token is #0
# Reduce with prod #0 [NT=0, SZ=2]
# Reduce rule: top state 0, lhs sym 0 -> state -1
# Goto state #-1

```





## Exercise

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Salad 2.10;  
Wine 12.00;  
Pasta 1.50;  
Bread 0.40;  
%  
Stefano : 2 Pasta, 1 Wine;  
Giulia : 1 Salad, 1 Bread, 1 Pasta;

/\* OUTPUT:  
Stefano: 15.0 EURO  
Giulia: 4.0 EURO  
\*/



## OTHER SLIDES



# Initializing parser  
 # Current Symbol is #2  
 # Shift under term #2 to state #1  
 # Current token is #3

----- ACTION\_TABLE -----  
 From state #0  
 [term 2:SHIFT(to state 1)]

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# Reduce with prod #3 [NT=2, SZ=1]  
 # Reduce rule: top state 0, lhs sym 2 -> state 2  
 # Goto state #2

----- ACTION\_TABLE -----  
 From state #1  
 [term 0:REDUCE(with prod 3)]  
 [term 3:REDUCE(with prod 3)]  
 ----- REDUCE\_TABLE -----  
 From state #0  
 [non term 1->state 3] [non term 2->state 2]

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