

POLITECNICO DI TORINO

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**(01JEUHT) Formal Languages and Compilers**  
Laboratory N°4


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Lab 4 1

## Attributes of Symbols

- A set of attributes can be associated to each symbol; attributes can be:
  - **Synthesized**: calculated from the values of the attributes of the node's children in the parse tree,
  - **Inherited**: calculated from the values of the parents / siblings in the parse tree.
- A set of semantic rules, specifying how attributes are calculated, is associated to each production.
- The scanner passes semantic values to the parser which, while recognizing the grammar, updates the nodes of the parse tree




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## Synthesized attributes

- A grammar whose attributes are all synthesized is denoted as an S-attribute grammar.
- In this case, it is possible to calculate the values of all attributes using a bottom-up strategy, from the leaves to the root of the parse tree.


E ::= E <sub>1</sub> '+' T	E.value = E <sub>1</sub> .value + T.value
E ::= T	E.value = T.value
T ::= number	T.value = number.value



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## Cup & Semantics: the Symbol class

- In Cup, each symbol in the stack is an object of class Symbol (cup/java\_cup/runtime/Symbol.java )
- It contains the following information:
  - A number uniquely identifying the symbol
    - ↳ public int sym;
  - The state in which the parse is
    - ↳ public int parse\_state;
  - Two integers that are used to pass the line and column number from the scanner to the parser
    - ↳ public int left, right;
  - An object of class Object to handle semantics
    - ↳ public Object value;



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scanner.jflex


## Passing semantic values to the parser

- Symbol and semantic value:  
[a-zA-Z][a-zA-Z0-9]\* { return new Symbol(sym.ID, new String(yytext())); }
- Symbol, line number, column number, and semantic value:  

```
%{
private Symbol my_symbol(int type, Object value){
return new Symbol(type, yyline, yycolumn, value);
}
%}
```

**Symbol Constructors:**  
 public Symbol(int sym\_id)  
 public Symbol(int sym\_id, int left, int right)  
 public Symbol(int sym\_id, Object o)  
 public Symbol(int sym\_id, int left, int right, Object o)

- Or equivalently:  
[a-zA-Z][a-zA-Z0-9]\* { return my\_symbol(sym.ID, new String(yytext())); }
- Or equivalently:  
[a-zA-Z][a-zA-Z0-9]\* { return new Symbol(sym.ID, yyline, yycolumn, new String(yytext())); }




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## Cup & Semantic: specifying nodes types

- Cup must know the type of the semantic value of each symbol
- It uses the following definition of terminals and non-terminals:
  - terminal <Object> <list\_of\_terminals> ;
  - non terminal <Object> <list\_of\_not\_terminals> ;
- <Object> is the class of the object associated to a given symbol
- Example:
  - terminal String ID;
    - ↳ An object of class String will be associated to ID.
  - terminal Integer NUM;
  - non terminal MyObject var;

```
class MyObject{
public String var_name;
public String var_type;
}
```



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### Cup & Semantic: using semantic values

- Given a set of productions:
 

```
E ::= E PLUS T
      | E MINUS T;
```
- One can refer to the semantic value of each symbol by adding labels to the symbols of interest:
  - A label is constituted by the ':' character followed by a name

```
E ::= E:n1 PLUS T:n2
      | E:n1 MINUS T:n2 ;
```
- Within each production, the labels can be used normally as objects of the class specified in the definition of terminals and non-terminals:
 

```
E ::= E:n1 PLUS T:n2 { : System.out.print(n1 + " + " + n2); ;
      | E:n1 MINUS T:n2 { : System.out.print(n1 + " - " + n2); ;
```

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### Cup & Semantic: Actions and RESULT

- An action can be associated to each production, ( { /\* Java Code\*/ ; } ) and is executed every time the corresponding production is reduced
- The action updates the semantic value of each symbol
- For each production, the RESULT object, of class Object, is defined.
- RESULT represents the result of the semantic rules contained in the action, **and is therefore associated to the symbol in the left hand side of the production**

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### Calculating synthesized attributes

- Given the algebraic expressions grammar, the following rule assigns to the symbol 'E' the sum or the subtraction of the values of the addends/subtrahends:
 

```
non terminal Integer E, T;
E ::= E:n1 PLUS T:n2
   { : RESULT = n1 + n2; ;
  | E:n1 MINUS T:n2
   { : RESULT = n1 - n2; ;
;
```
- OR: { : RESULT = new Integer(n1.intValue() + n2.intValue()); ;

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### Calculating synthesized attributes (2)

- It is possible to propagate more than one semantic value through RESULT, in the following way:
 

```
terminal RO, RC;
terminal String identifier;
terminal Integer Args;
non terminal Object[] Func;
non terminal goal;

goal ::= Func:a {
  System.out.println("Function name: " + a[0] + "Number of parameters: " + a[1]);
};

Func ::= identifier:a RO Args:b RC {
  RESULT = new Object[2];
  RESULT[0] = new String(a);
  RESULT[1] = new Integer(b);
};
```

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### Calculating synthesized attributes (3)

- Alternatively, one can write a class that contains all the required information:
 

```
action code {
  class MyFunc {
    public String id;
    public Integer args;
    MyFunc(String id, Integer args) {
      this.id = new String(id);
      this.args = new Integer(args);
    }
  }
};

non terminal MyFunc Func;

goal ::= Func:a {
  System.out.println("Function name : " + a.id + "Number of parameters: " + a.args );
};

Func ::= identifier:a RO Args:b RC { RESULT = new MyFunc( a, b ); ;
```

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### Parser debugging

- A series of options are available in Cup to visualize the parser's internal structures:
  - dump\_grammar : Prints the list of terminals, non-terminals and productions
  - dump\_states : Prints the state graph
  - dump\_table : Prints the ACTION TABLE and the REDUCE TABLE
  - dump : Prints all information
- The parser can be executed in debug mode (all the actions performed to analyze the input sequence are printed)

<b>Normal mode:</b> Yylex l = new Yylex(new FileReader(file)); parser p = new parser(l); Object result = p.parse();	<b>Debug mode:</b> Yylex l = new Yylex(new FileReader(file)); parser p = new parser(l); Object result = <b>p.debug_parse()</b> ;
--	---

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### States

**• -dump\_states**

```

===== Viable Prefix Recognizer =====
START lair_state [0]: {
[exp ::= (*) T, {EOF PLUS }]
[exp ::= (*) exp PLUS T, {EOF PLUS
}]
[T ::= (*) NUMBER, {EOF PLUS }]
[SSTART ::= (*) exp EOF, {EOF }]
}
transition on exp to state [3]
transition on T to state [2]
transition on NUMBER to state [1]
-----
lair_state [1]: {
[T ::= NUMBER (*), {EOF PLUS }]
}
-----
lair_state [2]: {
[exp ::= T (*), {EOF PLUS }]
}
-----
lair_state [3]: {
[exp ::= exp (*) PLUS T, {EOF PLUS
}]
[SSTART ::= exp (*) EOF, {EOF }]
}
transition on EOF to state [5]
transition on PLUS to state [4]
-----
lair_state [4]: {
[exp ::= exp PLUS (*) T, {EOF PLUS
}]
[T ::= (*) NUMBER, {EOF PLUS }]
}
transition on T to state [6]
transition on NUMBER to state [1]
-----
lair_state [5]: {
[SSTART ::= exp EOF (*), {EOF }]
}
-----
lair_state [6]: {
[exp ::= exp PLUS T (*), {EOF PLUS
}]
}

```

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### Action / Reduce Tables

**• -dump\_tables**

<pre> ----- ACTION_TABLE ----- From state #0 [term 2:SHIFT(to state 1)] From state #1 [term 0:REDUCE(with prod 3)] [term 3:REDUCE(with prod 3)] From state #2 [term 0:REDUCE(with prod 2)] [term 3:REDUCE(with prod 2)] From state #3 [term 0:SHIFT(to state 5)] [term 3:SHIFT(to state 4)] From state #4 [term 2:SHIFT(to state 1)] From state #5 [term 0:REDUCE(with prod 0)] From state #6 [term 0:REDUCE(with prod 1)] [term 3:REDUCE(with prod 1)] </pre>	<pre> ----- REDUCE_TABLE ----- From state #0 [non term 1-&gt;state 3] [non term 2-&gt;state 2] From state #1 From state #2 From state #3 From state #4 [non term 2-&gt;state 6] From state #5 From state #6 </pre>
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### Grammar

**• -dump\_grammar**

```

===== Terminals =====
[0]EOF [1]error [2]NUMBER [3]PLUS
===== Non terminals =====
[0]$START [1]exp [2]T
===== Productions =====
[0] $START ::= exp EOF
[1] exp ::= exp PLUS T
[2] exp ::= T
[3] T ::= NUMBER

```

```

exp → exp PLUS T
exp → T
T → NUMBER

```

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### Debugging

**• debug\_parse()**      **Input string: 3+5**

<pre> # Initializing parser FOUND: 3 # Current Symbol is #2 # Shift under term #2 to state #1 FOUND: + # Current token is #3 # Reduce with prod #3 [NT=2, SZ=1] # Reduce rule: top state 0, lhs sym 2 -&gt; state 2 # Goto state #2 # Reduce with prod #2 [NT=1, SZ=1] # Reduce rule: top state 0, lhs sym 1 -&gt; state 3 # Goto state #3 # Shift under term #3 to state #4 FOUND: 5 # Current token is #2 </pre>	<pre> # Shift under term #2 to state #1 # Current token is #0 # Reduce with prod #3 [NT=2, SZ=1] # Reduce rule: top state 4, lhs sym 2 -&gt; state 6 # Goto state #6 Found expression # Reduce with prod #1 [NT=1, SZ=3] # Reduce rule: top state 0, lhs sym 1 -&gt; state 3 # Goto state #3 ----- # Shift under term #0 to state #5 # Current token is #0 # Reduce with prod #0 [NT=0, SZ=2] # Reduce rule: top state 0, lhs sym 0 -&gt; state -1 # Goto state #-1 </pre>
--	--

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### Exercise

```

Salad 2.10;
Wine 12.00;
Pasta 1.50;
Bread 0.40;
%
Stefano : 2 Pasta, 1 Wine;
Giulia : 1 Salad, 1 Bread, 1 Pasta;

/* OUTPUT:
Stefano: 15.0 EURO
Giulia: 4.0 EURO
*/

```

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### OTHER SLIDES

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